

# Residuated frames for substructural logics, Part III

**Peter Jipsen**, Chapman University

joint work with **N. Galatos**, University of Denver, Colorado,  
**M. A. Moshier**, Chapman University, California and  
**G. Greco, F. Liang, A. Palmigiano, A. Tzimoulis**,  
Delft University of Technology, the Netherlands

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# Overview

## Part III

- Introduction to algebraic proof theory
- Okada lemma
- Decidability of equational theory of RL
- Finite model property
- Finite embeddability property

## Part IV

- Categorical duality for complete lattices
- Duality for join-semilattices and algebraic frames
- Correspondence theory briefly
- De Morgan lattices, ortholattices and involutive frames
- Involutive FL-algebras and involutive residuated frames
- Bunched implication algebras and their frames
- Open problems and projects

# The sequent calculus **GN**

$$\begin{array}{c}
 \frac{xNa \quad aNz}{xNz} \text{ (CUT)} \quad \frac{}{aNa} \text{ (Id)} \\
 \\
 \frac{xNa \quad bNz}{x \circ (a \setminus b)Nz} (\setminus L) \quad \frac{a \circ xNb}{xNa \setminus b} (\setminus R) \\
 \\
 \frac{xNa \quad bNz}{(b/a) \circ xNz} (/L) \quad \frac{x \circ aNb}{xNb/a} (/R) \\
 \\
 \frac{a \circ bNz}{a \cdot bNz} (\cdot L) \quad \frac{xNa \quad yNb}{x \circ yNa \cdot b} (\cdot R) \\
 \\
 \frac{aNz}{a \wedge bNz} (\wedge L\ell) \quad \frac{bNz}{a \wedge bNz} (\wedge Lr) \quad \frac{xNa \quad xNb}{xNa \wedge b} (\wedge R) \\
 \\
 \frac{aNz \quad bNz}{a \vee bNz} (\vee L) \quad \frac{xNa}{xNa \vee b} (\vee R\ell) \quad \frac{xNb}{xNa \vee b} (\vee Rr) \\
 \\
 \frac{\varepsilon Nz}{1Nz} (1L) \quad \frac{}{\varepsilon N1} (1R)
 \end{array}$$

$a, b \in Fm$  or  $B$ ,  $x, y \in W$ ,  $z \in W'$ ,  $xNa$  is short for  $xN(id, a)$  meaning  $x \Rightarrow a$ .

## Semantics of sequent calculi: Residuated frames

Let  $\mathbf{GN}_{cf}$  be the sequent calculus **GN** without the **cut rule**

Define a binary relation  $N \subseteq W \times W' = F_{\text{Mon}(\circ, \varepsilon)}(Fm) \times (S_W \times Fm)$  by

$$wN(u, a) \iff u[w] \Rightarrow a \quad \text{is provable in } \mathbf{GN}_{cf}$$

Define the **accessibility** relations  $\circ \subseteq W^3, \backslash, /$  by

$$\circ(v_1, v_2, w) \iff v_1 \circ v_2 = w$$

$$\backslash = \{((u, a), x, (u[\_ \circ x], a)) : u \in W, a \in Fm, x \in W\}$$

$$/ = \{(x, (u, a), (u(x \circ \_), a)) : u \in W, a \in Fm, x \in W\}$$

$$E = \{\varepsilon\}$$

**Lemma 36:**  $\mathbf{W}_{RL} = (W, W', N, \circ, \backslash, /, E)$  is a **residuated frame**

# Algebraic cut-admissibility

Theorem (Okada, Terui 1999, Galatos, J. 2013)

The following are equivalent:

- 1  $t \Rightarrow a$  is provable in **GN**
- 2  $t \leq a$  holds in **RL**
- 3  $t \Rightarrow a$  is provable in **GN<sub>cf</sub>**

**Proof** (outline): (3 $\Rightarrow$ 1) is obvious. (1 $\Rightarrow$ 2) Assume  $t \Rightarrow a$  is provable **with cut**. Show that **all sequent rules** hold as quasiequations in **RL** (where  $\Rightarrow, \circ$  are **replaced by**  $\leq, \cdot$ )

(2 $\Rightarrow$ 3) Assume  $t \leq a$  holds in **RL** and define the algebra

$\mathbf{W}_{\mathbf{RL}}^+ = (\gamma_N[\mathcal{P}(W)], \cup_{\gamma_N}, \cap, \cdot, 1, \setminus, /)$  using the **closed sets**  $\gamma_N(X)$  of the **polarity frame**  $(W, W', N)$  and  $1 = \gamma_N(\{\varepsilon\})$

$X \cdot Y = \gamma_N(\{w : \circ(u, v, w) \text{ for some } u \in X, v \in Y\})$

$X \setminus Y = \{w \in W : X \cdot \{w\} \subseteq Y\}$        $Y / X = \{w \in W : \{w\} \cdot X \subseteq Y\}$ .

## Proof outline (continued)

Then  $\mathbf{W}_{\mathbf{RL}}^+$  is a residuated lattice, hence satisfies  $t \leq a$ .

Let  $f : Fm \rightarrow \mathbf{B} = Fm$  be any **homomorphism**.

**Extend** to  $\bar{f} : W \rightarrow \mathbf{W}_{\mathbf{RL}}^+$  by  $\bar{f}(p) = N^\downarrow\{f(p)\}$ , so  $t \leq a$  implies  $\bar{f}(t) \subseteq \bar{f}(a)$

Recall, for all  $b \in Fm$ ,  $N^\downarrow\{b\} = \{w \in W : wN(x_0, b)\}$ .

Prove by **induction** that  $f(b) \in \bar{f}(b) \subseteq N^\downarrow\{f(b)\}$  for all  $b \in Fm$ .

Then  $f(t) \in \bar{f}(t) \subseteq \bar{f}(a) \subseteq N^\downarrow\{f(a)\}$ , hence  $f(t)N(x_0, f(a))$ .

Using  $f = id$ , it follows that  $t \Rightarrow a$  holds in  $\mathbf{GN}_{cf}$ . □

## Gentzen frames

The following properties hold for  $\mathbf{W}_L$ ,  $\mathbf{W}_{RL}$  and  $\mathbf{W}_{A,B}$ :

$\mathbf{W}$  is a residuated frame

$\mathbf{B}$  is a (partial) algebra of the same type, ( $\mathbf{B} = \mathbf{L}, \mathbf{Fm}, \mathbf{B}$ )

$B$  generates  $(W, \circ, \varepsilon)$  (as a monoid)

$W'$  contains a copy of  $B$  ( $b \leftrightarrow (id, b)$ )

$N$  satisfies **GN**, for all  $a, b \in B$ ,  $x, y \in W$ ,  $z \in W'$ .

We call such pairs  $(\mathbf{W}, \mathbf{B})$  *Gentzen frames*.

A *cut-free Gentzen frame* is not assumed to satisfy the (CUT)-rule.

**Theorem.** Given a Gentzen frame  $(\mathbf{W}, \mathbf{B})$ , the map  $N^\downarrow\{\cdot\} : \mathbf{B} \rightarrow \mathbf{W}^+$ ,  $b \mapsto N^\downarrow\{b\}$  is a (partial) homomorphism.

(Namely, if  $a, b \in B$  and  $a \bullet b \in B$  ( $\bullet$  is a connective) then  $N^\downarrow\{a \bullet_{\mathbf{B}} b\} = N^\downarrow\{a\} \bullet_{\mathbf{W}^+} N^\downarrow\{b\}$ ).

## Proof

**Key Lemma.** Let  $(\mathbf{W}, \mathbf{B})$  be a cut-free Gentzen frame. For all  $a, b \in B$ ,  $X, Y \in \mathbf{W}^+$  and for every connective  $\bullet$ , if  $a \bullet b \in B$ ,  $a \in X \subseteq N^\downarrow\{a\}$  and  $b \in Y \subseteq N^\downarrow\{b\}$ , then

$$a \bullet_{\mathbf{B}} b \in X \bullet_{\mathbf{W}^+} Y \subseteq N^\downarrow\{a \bullet_{\mathbf{B}} b\} \text{ (and } 1_{\mathbf{B}} \in 1_{\mathbf{W}^+} \subseteq N^\downarrow\{1_{\mathbf{B}}\} \text{ )}$$

In particular,  $a \bullet_{\mathbf{B}} b \in N^\downarrow\{a\} \bullet_{\mathbf{W}^+} N^\downarrow\{b\} \subseteq N^\downarrow\{a \bullet_{\mathbf{B}} b\}$ .

Furthermore, given (CUT), we have equality.

**Proof** Let  $\bullet = \vee$ . If  $x \in X$ , then  $x \in N^\downarrow\{a\}$ ; so  $xNa$  and  $xNa \vee b$ , by ( $\vee R$ ); hence  $x \in N^\downarrow\{a \vee b\}$  and  $X \subseteq N^\downarrow\{a \vee b\}$ . Likewise  $Y \subseteq N^\downarrow\{a \vee b\}$ , so  $X \cup Y \subseteq N^\downarrow\{a \vee b\}$  and  $X \vee Y = \gamma(X \cup Y) \subseteq N^\downarrow\{a \vee b\}$ .

On the other hand, let  $X \vee Y \subseteq N^\downarrow\{z\}$ , for some  $z \in W$ . Then,  $a \in X \subseteq X \vee Y \subseteq N^\downarrow\{z\}$ , so  $a Nz$ . Similarly,  $b Nz$ , so  $a \vee b Nz$  by ( $\vee L$ ), hence  $a \vee b \in N^\downarrow\{z\}$ . Thus,  $a \vee b \in X \vee Y$ .

We used that every closed set is an intersection of *basic closed sets*  $N^\downarrow\{z\}$ , for  $z \in W$ .



# Applications of frames: MacNeille-completion

For a **residuated lattice**  $\mathbf{L}$ , we associated the Gentzen frame  $(\mathbf{W}_{\mathbf{L}}, \mathbf{L})$ .

The underlying poset of  $\mathbf{W}_{\mathbf{L}}^+$  is the *MacNeille completion* of the underlying poset reduct of  $\mathbf{L}$ .

**Theorem.** The map  $x \mapsto N^{\downarrow}\{x\}$  is an embedding of  $\mathbf{L}$  into  $\mathbf{W}_{\mathbf{L}}^+$ .

(A general residuated frame is )

## Completeness - Cut elimination

**Corollary.**  $\mathbf{RLsq}$  is complete with respect to  $\mathbf{W}_{\mathbf{RL}}^+$ .

**Corollary.** The algebra  $\mathbf{W}_{\mathbf{RL}}^+$  generates the variety  $\mathbf{RL}$ .

The frame  $\mathbf{W}_{\mathbf{RL}_{cf}}$  corresponds to cut-free  $\mathbf{RLsq}_{cf}$ .

**Corollary (Cut Elimination).**  $\mathbf{RLsq}$  and  $\mathbf{RLsq}_{cf}$  prove the same sequents.

**Corollary.** The equational theory of  $\mathbf{RL}$  is decidable.

## Finite model property, FMP

For  $\mathbf{W}_{\mathbf{RLsq}}$ , given  $(x, z) \in W \times W'$  (if  $z = (u, c)$ , then  $u(x) \Rightarrow c$  is a sequent), we define  $(x, z)^\uparrow$  as the smallest subset of  $W \times W'$  that contains  $(x, z)$  and is closed upwards with respect to the rules of  $\mathbf{RLsq}_{cf}$ . Note that  $(x, z)^\uparrow$  is finite.

A new frame  $\mathbf{W}^\uparrow$  associated with  $N' = N \cup ((x, z)^\uparrow)^c$  is residuated and Gentzen.

$(N')^c$  is finite, so has **finite** domain  $Dom((N')^c)$  and codomain  $Cod((N')^c)$

For every  $y \notin Cod((N')^c)$ ,  $N'^\downarrow\{y\} = W$ . So,  $\{N'^\downarrow\{y\} : y \in W\}$  is finite and a basis for  $\gamma_{N'}$ . So  $\mathbf{W}^{\uparrow+}$  is finite.

Moreover, if  $u(x) \Rightarrow c$  is not provable in  $\mathbf{RLsq}$ , then it is not valid in  $\mathbf{W}^{\uparrow+}$ .

**Corollary** . The system  $\mathbf{RLsq}$  has the **finite model property**.

**Corollary**. The variety  $\mathbf{RL}$  is generated by its **finite members**.

## Finite embedding property, FEP

A class of algebras  $\mathcal{K}$  has the **finite embeddability property (FEP)** if for every  $\mathbf{A} \in \mathcal{K}$ , every finite partial subalgebra  $\mathbf{B}$  of  $\mathbf{A}$  can be (partially) embedded in a finite  $\mathbf{D} \in \mathcal{K}$ .

The corresponding logic has the *strong finite model property*:

if  $\Phi \not\vdash \psi$ , for finite  $\Phi$ , then there is a finite counter-model, namely there is  $\mathbf{D} \in \mathcal{K}$  and a homomorphism  $f : \mathbf{Fm} \rightarrow \mathbf{D}$ , such that  $f(\phi) = 1$ , for all  $\phi \in \Phi$ , but  $f(\psi) \neq 1$ .

For logics with finitely many axioms and rules, this implies that the *deducibility relation* is decidable

For a finitely axiomatized variety, FEP implies the decidability of the *universal theory*

## FEP for integral RLs with $\{\vee, \cdot, 1\}$ -equations

Blok and van Alten 2002 proved FEP for integral RLs, and extended it to residuated groupoids (2005)

**Theorem.** Every variety of *integral RL's* axiomatized by equations over  $\{\vee, \cdot, 1\}$  has the FEP.

$\mathbf{B}$  embeds in  $\mathbf{W}_{\mathbf{A},\mathbf{B}}^+$  via  $N^\downarrow\{\_ \} : \mathbf{B} \rightarrow \mathbf{W}_{\mathbf{A},\mathbf{B}}^+$

$\mathbf{W}_{\mathbf{A},\mathbf{B}}^+$  is finite

$\mathbf{W}_{\mathbf{A},\mathbf{B}}^+ \in \mathcal{V}$

**Corollary.** These varieties are generated as quasivarieties by their finite members.

**Corollary.** The corresponding logics have the *strong finite model property*

# Finiteness

Idea for finiteness: Every element in  $\mathbf{W}_{A,B}^+$  is an intersection of basic elements. So it suffices to prove that there are only finitely many such elements.

Replace the frame  $\mathbf{W}_{A,B}$  by one  $\mathbf{W}_{A,B}^M$ , where it is easier to work.

Let  $\mathbf{M}$  be the free monoid with unit over the set  $B$  and  $f : M \rightarrow W$  the extension of the identity map.

$$M \xrightarrow{f} W \xrightarrow{N} W'.$$

$\mathbf{M}$  is partially ordered by  $s \leq^{\mathbf{M}} t$  if  $s$  is a subsequence of  $t$ .

Blok and van Alten show that this order has no infinite antichains and no infinite ascending chains, so by Higman's lemma it is dually well-ordered.

Integrality ( $x \leq 1$ ) is used to show that any downset in  $\mathbf{M}$  is  $\downarrow M_0$  for some finite set  $M_0 \subseteq M$ .

# Equations 1

Idea: Express equations over  $\{\vee, \cdot, 1\}$  at the frame level.

For an equation  $\varepsilon$  over  $\{\vee, \cdot, 1\}$  we distribute products over joins to get  $s_1 \vee \cdots \vee s_m = t_1 \vee \cdots \vee t_n$ .  $s_i, t_j$ : monoid terms.

$s_1 \vee \cdots \vee s_m \leq t_1 \vee \cdots \vee t_n$  and  $t_1 \vee \cdots \vee t_n \leq s_1 \vee \cdots \vee s_m$ .

The first is equivalent to:  $\&(s_j \leq t_1 \vee \cdots \vee t_n)$ .

We proceed by example:  $x^2y \leq xy \vee yx$

$$(x_1 \vee x_2)^2y \leq (x_1 \vee x_2)y \vee y(x_1 \vee x_2)$$

$$x_1^2y \vee x_1x_2y \vee x_2x_1y \vee x_2^2y \leq x_1y \vee x_2y \vee yx_1 \vee yx_2$$

$$x_1x_2y \leq x_1y \vee x_2y \vee yx_1 \vee yx_2$$

$$\frac{x_1y \leq v \quad x_2y \leq v \quad yx_1 \leq v \quad yx_2 \leq v}{x_1x_2y \leq v}$$

$$\frac{x_1 \circ y \ N \ z \quad x_2 \circ y \ N \ z \quad y \circ x_1 \ N \ z \quad y \circ x_2 \ N \ z}{x_1 \circ x_2 \circ y \ N \ z} \quad R(\varepsilon)$$

## Equations 2

**Theorem.** If  $(\mathbf{W}, \mathbf{B})$  is a Gentzen frame and  $\varepsilon$  an equation over  $\{\vee, \cdot, 1\}$ , then  $(\mathbf{W}, \mathbf{B})$  satisfies  $R(\varepsilon)$  iff  $\mathbf{W}^+$  satisfies  $\varepsilon$ .

(The linearity of the denominator of  $R(\varepsilon)$  plays an important role in the proof.)

**Corollary.** If an equation over  $\{\vee, \cdot, 1\}$  is valid in  $\mathbf{A}$ , then it is also valid in  $\mathbf{W}_{\mathbf{A}, \mathbf{B}}^+$ , for every partial subalgebra  $\mathbf{B}$  of  $\mathbf{A}$ .

Consequently,  $\mathbf{W}_{\mathbf{A}, \mathbf{B}}^+ \in \mathcal{V}$ .



## Structural rules

Given an equation  $\varepsilon$  of the form  $t_0 \leq t_1 \vee \dots \vee t_n$ , where  $t_i$  are  $\{\cdot, 1\}$ -terms we construct the rule  $R(\varepsilon)$

$$\frac{u[t_1] \Rightarrow a}{u[t_0] \Rightarrow a} (R(\varepsilon))$$

where the  $t_i$ 's are evaluated in  $(W, \circ, \varepsilon)$ . Such a rule is called *linear* if all variables in  $t_0$  are distinct.

**Theorem.** Every system obtained from **RLsq** by adding linear rules has the cut elimination property.

A set of rules of the form  $R(\varepsilon)$  is called *reducing* if there is a complexity measure that decreases with upward applications of the rules (and the rules of **RLsq**).

**Theorem.** Every system obtained from **RLsq** by adding linear reducing rules is decidable. The subvariety of residuated lattices axiomatized by the corresponding equations has decidable equational theory.

# Applications

- **Cut-elimination** (CE) and **finite model property** (FMP) for **FL** and (cyclic) **InFL**. Generation by finite members for RL, InFL
- M. Kozak 2008 proved **distributive FL** has the FMP, and using our approach the same result holds for any extension of **DFL** with linear reducing structural rules
- The **finite embeddability property** (FEP) for integral RL with  $\{\vee, \cdot, 1\}$ -axioms
- The above extend to the **non-associative case**, also with the addition of suitable **structural rules**

## Recent papers using residuated frames

**Terui** 2018 MacNeille completion and Buchholz' Omega rule for parameter-free second order logics, arXiv:1804.11066v1

**Galatos and Prenosil** 2018 On an equivalence between integral and involutive residuated structures, preprint

**Buszkowski** 2017 Involutive nonassociative Lambek calculus, sequent systems and complexity, Bull. Sec. of Logic, 46, 75–91

**Cardona and Galatos** 2017 The FEP for some varieties of fully distributive knotted residuated lattices, Alg. Univ., 78, 363–376

**Ciabattoni, Galatos and Terui** 2017 Algebraic proof theory: hypersequents and hypercompletions, Ann. of Pure and App. Logic, 168

**Baldi and Terui** 2016 Densification of FL chains via residuated frames, Alg. Univ., 75(2), 169–195

**Horčík** 2015 Word problem for knotted residuated lattices, J. of Pure and App. Alg., 219, 1548–1563

**Ciabattoni, Galatos and Terui** 2012 Algebraic proof theory for substructural logics: cut-elimination and completions, Ann. of Pure and App. Logic, 163, 266–290

## Filter-ideal frames and canonical extensions

Let  $L$  be a (bounded) lattice.

Recall that the **Dedekind-MacNeille frame** is  $\mathbf{W}_L = (L, L, \leq)$ .

$\mathbf{W}_L^+$  is the **MacNeille completion** of  $L$ .

A **filter**  $F \subseteq L$  is a meet-closed upset:  $z \geq x, y \in F \implies x \wedge y, z \in F$ .

**Ideals** of  $L$  are defined dually.

Let  $\mathcal{F}L = \{\text{nonempty filters of } L\}$ ,  $\mathcal{I}L = \{\text{nonempty ideals of } L\}$

Define  $N \subseteq \mathcal{F}L \times \mathcal{I}L$  by  $FND$  iff  $F \cap D \neq \emptyset$

The **filter-ideal frame** is  $FI_L = (\mathcal{F}L, \mathcal{I}L, N)$

The **canonical extension** of  $L$  is  $L^\delta = FI_L^+$

## Completions

A **completion** of a lattice  $L$  is an embedding  $e : L \rightarrow \bar{L}$  where  $\bar{L}$  is a complete lattice

The MacNeille completion is **join-dense** and **meet-dense**  
(every  $x \in \mathbf{W}_L^+$  is a **join** of elements of  $L$  and a **meet** of elements of  $L$ )

$\implies$  all existing joins and meets of  $L$  are **preserved**

The **canonical extension** is **dense** (every  $x \in L^\delta$  is a **join of a meet** of elements of  $L$  and a **meet of a join** of elements of  $L$ )

and **compact** (if  $\bigwedge A \leq \bigvee B$  for  $A, B \subseteq L$  then  $a_1 \wedge \dots \wedge a_m \leq b_1 \vee \dots \vee b_m$  for some  $a_i \in A, b_j \in B$ ).

$\implies$  all infinite joins and meets of  $L$  are destroyed

So these polarity frames are **two extremes** of lattice completions

Polarity frames in general are a good way to study **dense completions**