

# The Duality of Time and Information in Concurrency and Branching Time

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# 1. Outline

- 1 Chu spaces
- 2 Applications to program semantics: schedules and their dual automata
- 3 Chu-like categories
- 4 Presheaf-like categories
- 5 Typed Chu spaces
- 6 Applications to philosophy: dualism, properties, qualia

## 2. Chu spaces: Motivation

- Duality. ( $\mathbf{Chu}(\mathbf{Set}, K)$  is a self-dual category).
- Generality. (Every category of relational structures of total arity  $n$  is embeds concretely in  $\mathbf{Chu}(\mathbf{Set}, 2^n)$ ).
- Real-world applications, as e.g. per the title.

### 3. Chu spaces: Definition

Given a set  $K$ , a **Chu space over  $K$**  is a triple  $(A, r, X)$  such that  $A$  and  $X$  are sets and  $r$  is a function  $r : A \times X \rightarrow K$  defining an  $A \times X$  matrix over  $K$ . That is,  $r(a, x)$  is the element at row  $a$  and column  $x$ .

*Duality.* The dual of a Chu space  $(A, r, X)$  is the Chu space  $(A, r, X) \perp = (X, r^\smile, A)$ .

## 4. Uncertainty principle for Chu spaces

A Chu space over  $2 = 0, 1$  cannot have:

- 1 an all-ones row and an all-zeroes column, and dually;
- 2 closure of rows under OR and columns under AND, and dually.

## 5. Extensionality

- Extensional: No duplicate rows
- Normal: No duplicate columns
- Biextensional: extensional and normal

## 6. Representation of lattice-like objects as biextensional Chu spaces

Set: A discrete Chu space (no omitted columns) s.t.  $|K| \geq 2$ .

CABAs: A codiscrete Chu space (no omitted rows).  $|K| \geq 2$ .

Remaining examples take  $K = \mathbf{2} = \{0, 1\}$ .

Poset: columns closed under arbitrary OR and AND.

Distributive lattice: rows closed under arbitrary OR and AND

Complete join-semilattice: rows and columns closed under arbitrary OR.

Complete meet-semilattice: rows and columns closed under arbitrary AND.

T0 topological space: columns closed under arbitrary OR and finite AND.

Finite-dimensional vector space over GF2: finite, rows and columns closed under XOR. Necessarily  $2^n \times 2^n$ . (Infinite-dimensional case need not be square.)

Many other examples.

## 7. Properties of Chu spaces

A property of a Chu space is a set of omitted columns.

The property *true* is the empty set.

The maximal property is the set of all omitted columns.

Remark. The properties of a Chu space form a CABA.

The only property of a set is *true* (the degenerate CABA).

The properties of any other structure form a nondegenerate CABA.



## 8. Chu transforms $(f, g) : (A, r, X) \rightarrow (B, s, Y)$

Biextensional case: An  $A \times Y$  matrix whose rows come from  $(B, s, Y)$  and whose columns come from  $(A, r, X)$ . Row  $a$  of  $(f, g)$  equals row  $f(a)$  of  $B$  while column  $y$  equals column  $g(y)$  of  $(A, r, X)$ .

Image of  $(f, g) : (f(A), s', Y)$  s.t.  $s'(b', y) = s(b, y)$  for all  $b$  in  $f(A)$  and all  $y$  in  $Y$ .

### Theorem

*Any property of a Chu space is a property of its image under a Chu transform.*

### Proof.

Any omitted column of the source must be omitted from the transform. Hence its image under  $f$  must be omitted from the subspace  $f(A)$  of the target. □

General case (no extensionality needed): A Chu transform is a pair  $(f, g)$  where  $f : A \rightarrow B, g : Y \rightarrow X$  s.t.  $s(f(a), y) = r(a, g(y))$  for all  $a$  in  $A, y$  in  $Y$  (adjointness).

## 9. Comonoids

A **comonoid** is a Chu space  $(A, r, X)$  over  $\mathbf{2}$  such that

- (i) (The counit) its states include the all-zeroes and all-ones columns.
- (ii) (The comultiplication, aka the crossword property) Given any  $A \times A$  matrix whose rows and columns are drawn from the columns of  $(A, r, X)$ , its main diagonal is a column of  $X$ .

A T1 comonoid is one s.t. the underlying poset of the rows is discrete.

## 10. Are there any hard problems about Chu spaces?

*Problem* (open for 20 years). Is every T1 comonoid discrete?

True for countable comonoids (known in 1995).

In 2015 Pace Nielsen answered this in the negative with a counterexample of cardinality  $\aleph_1$ . Accepted by JLMS.

George M. Bergman, Pace P. Nielsen, *On Vaughan Pratt's crossword problem*, Journal of the London Mathematical Society, Volume 93, Issue 3, 1 June 2016, Pages 825845, <https://doi.org/10.1112/jlms/jdw011>.

# 11. Interpretation as schedules and their dual automata

Simplest case:  $K = \mathbf{2} = \{0, 1\}$ .

Define a schedule to be a poset, and an automaton to be a distributive lattice. Interpret the rows of a schedule as its events, with the order constraining their order in time. Interpret its columns as the possible states of its dual automaton.

## 12. True concurrency via transition $T$

$$K = \{0, T, 1\} = \mathbf{3}.$$

Motivation: refine before-after to before-during-after,  $T$  being the state of transition of an event.

A connected event is the discrete  $1 \times 3$  Chu space  $0T1$  over  $\mathbf{3}$ , having three states, before, during, and after. Here  $T$  is understood geometrically as an edge connecting  $0$  and  $1$ .

Represent true concurrence of two events as the discrete two-point Chu space over  $\mathbf{3}$ . Its dual automaton,  $9 \times 2$ , consists of the 4 quiescent states  $00, 01, 10, 11$ , the 4 semi-quiescent states  $0T, T0, 1T, T1$ , and the truly concurrent state  $TT$ .

Represent *mutex*( $a, b$ ) (mutual exclusion) as the  $2 \times 8$  Chu space over  $\mathbf{3}$  whose dual  $8 \times 2$  automaton omits the row  $TT$ .

Each schedule over  $2$  gives rise to many schedules over  $\mathbf{3}$  according which states containing  $T$  are omitted.

This gives a denotational semantics for higher dimensional automata, with the dimension of a state being simply the number of  $T$ 's in it.

## 13. Branching time via cancellation

Motivation: Prior denotational semantics of branching time were all higher order and obscure. before-during-after-cancelled, as in “the event has been cancelled”, represented by state  $X$ .

Distinguish  $ab + ac$  from  $a(b + c)$  as having the additional property that one of  $b$  or  $c$  must be cancelled before  $a$  starts.

## 14. Set-like categories

Motivation: A new way of defining  $\mathbf{Chu}(\mathbf{Set}, K)$ .

A locally small category  $\mathcal{C}$  is **set-like** when it contains an object  $\mathbf{1}$  such that

(i)  $|\mathcal{C}(\mathbf{1}, \mathbf{1})| = 1$  ( $\mathbf{1}$  is rigid).

(ii) (Extensionality) For any two morphisms  $f, g : X \rightarrow Y$  s.t.  $fx = gx$  for all  $x : \mathbf{1} \rightarrow X$ ,  $f = g$ .

Interpretation: Each object  $X$  of  $\mathcal{C}$  represents the set  $\mathcal{C}(\mathbf{1}, X)$  of elements of  $X$ , while each morphism  $f : X \rightarrow Y$  represents the function defined by its left action on the elements of  $X$ .

Take **Set** to be the maximal set-like category whose objects are all homsets of all locally small categories. (Maximality fills up the homsets. Circular, but all definitions of **Set** are inevitably circular in one way or another.)

## 15. Chu-like categories

A locally small category  $\mathcal{C}$  is **Chu-like over**  $K$  when it contains objects  $\mathbf{1} = \mathbf{1}!K$  and  $\perp = K!1$  such that

- (i)  $\mathbf{1}$  and  $\perp$  are rigid.
- (ii) (Biextensionality) For any two morphisms  $f, g : ArX \rightarrow BsY$ , if  $yfa = yga$  for all  $a : \mathbf{1} \rightarrow ArX$  and all  $y : BsY \rightarrow \perp$ , then  $f = g$ .

Interpret  $K$  as  $\mathcal{C}(\mathbf{1}, \perp)$ , For any object  $ArX$ , interpret  $\mathcal{C}(\mathbf{1}, ArX)$  as  $A$ ,  $\mathcal{C}(ArX, \perp)$  as  $X$ , and  $xa$  as  $r(a, x)$  for all  $a \in A, x \in X$ .

Interpret every homomorphism  $h : ArX \rightarrow BsY$  as the pair  $(f, g)$  where  $f : A \rightarrow B$  is the left action of  $h$  on  $\mathcal{C}(\mathbf{1}, ArX)$  and  $g$  is the right action of  $h$  on  $\mathcal{C}(BSY, \perp)$ . That is,  $f(a) = ha$  and  $g(y) = yh$ .

Every homomorphism  $h : ArX \rightarrow BsY$  satisfies adjointness as a consequence of associativity of composition, as displayed thus.

$$\mathbf{1} \xrightarrow{a} ArX \xrightarrow{h} BsY \xrightarrow{y} \perp$$

The Chu-like categories over  $K$  are the subcategories  $\mathcal{C}$  of  $\mathbf{Chu}(\mathit{Set}, K)$  defined analogously to the set-like categories.



## 16. Grph-like categories

A locally small category  $\mathcal{C}$  is **Grph-like** when it contains objects  $V, E$  such that

(i)  $V$  and  $E$  are rigid;

(ii)  $\mathcal{C}(V, E) = \{s, t\}$  and  $\mathcal{C}(E, V)$  is empty;

(iii) (Extensionality) For any two morphisms  $f, g : G \rightarrow H$ , if  $fv = gv$  for all  $v : V \rightarrow G$ , and  $fe = ge$  for all  $e : E \rightarrow G$ , then  $f = g$ .

The **Grph-like** categories are the subcategories  $\mathcal{C}$  of **Grph** that contain a graph  $V$  consisting of one vertex  $v$  and no edges, and a graph  $E$  consisting of one edge  $e$  and distinct vertices  $es, et$ , such that the subcategory is full on the carriers  $\mathcal{C}(V, G)$  (the vertices of  $G$ ) and  $\mathcal{C}(E, G)$  (the edges of  $G$ ).

## 17. Presheaf-like categories

These generalize **Grph**-like categories uniformly.

Given a small category  $\mathcal{J}$ , a presheaf-like category  $\mathcal{C}$  is one for which  $\mathcal{J}$  is a full subcategory of  $\mathcal{C}$  and the morphisms of  $\mathcal{C}$  are extensional on homsets  $\mathcal{C}(j, X)$  for all  $j \in \text{ob}(\mathcal{J})$ .

## 18. Typed Chu spaces

These are a common generalization of Chu-like categories and presheaf-like categories.

Given a profunctor (distributor, (bi)module)  $K : \mathcal{L} \nrightarrow \mathcal{J}$  between two small categories  $\mathcal{L}$  and  $\mathcal{J}$ , a typed-Chu-space-like category  $\mathcal{C}$  is one for which  $K$  is a subcategory of  $\mathcal{C}$  whose objects are those of  $L = \text{ob}(\mathcal{L})$  and  $J = \text{ob}(\mathcal{J})$  and whose morphisms are all morphisms  $k : J \rightarrow L$  of  $K$ . (The notation  $\mathcal{L} \nrightarrow \mathcal{J}$  is per convention, though Peter Johnstone prefers  $\mathcal{J} \nrightarrow \mathcal{L}$ , reasonably.)

This creates a framework in which the objects of  $\mathcal{J}$  serve as the sorts of the points of a typed Chu space while the objects of  $\mathcal{L}$  serve as the properties of its states. The adjointness chain generalizes as follows.

$$j' \xrightarrow{f} j \xrightarrow{a} A \xrightarrow{h} B \xrightarrow{y} \ell \xrightarrow{\pi} \ell'$$

## 19. Applications to philosophy

- 1 Cartesian dualism:  $A$  as body,  $X$  as mind,  $r : A \times X \rightarrow K$  as a mathematical pituitary gland.
- 2 (Co)extensionality of properties via their *possible* states.
- 3 Qualia as the morphisms  $k : j \rightarrow \ell$ , e.g. the possible colors of a cat, the possible heights of buildings, . . . . Qualia are both points and states, expressing C.I.Lewis's intuitions about the psychological vs. physical ambiguity of the nature of qualia.

THANK YOU